

## II. Public Key Cryptography Algorithms

• Uses pair of keys:

- Public Key for encryption ( $K_U$ )
- Private Key for decryption ( $K_R$ )

• RSA (Rivest - Shamir - Adleman)

→ Key Generation

- 1) Select two large prime numbers  $p$  and  $q$
- 2) Compute  $n = p \times q$
- 3) Compute Euler's function  $\phi(n) = (p-1)(q-1)$
- 4) Choose a public exponent  $e$  such that  $1 < e < \phi(n)$  and  $\gcd(e, \phi(n)) = 1$
- 5) Compute the private key  $d$  such that
$$d \times e \equiv 1 \pmod{\phi(n)}$$
$$d \times e = 1 + k \times \phi(n)$$
- 6)  $K_U: \{e, n\}$  ,  $K_R: \{d, n\}$

Notes:

- $d$  should be : integer ,  $< \phi(n)$
- $k$  is just an integer (e.g., 1, 2, 3, 4, ...)  
that helps us solve  $d$
- We try different options of  $k$  to find a valid solution
- $e$  and  $d$  should not be the same

→ Encryption:  $C = M^e \pmod{n}$

→ Decryption:  $M = C^d \pmod{n}$

Note:

•  $M$  is converted to a nb. based on the alphabetical order

6) Public Key:  $(e, n)$  , Private Key:  $(d, n)$

RSA Encryption and Decryption

→ Encryption (Sender Side)

$$C = M^e \pmod n$$

M: Message (as a nb.)

C: Encryption message

→ Decryption (Receiver Side)

$$M = C^d \pmod n$$

Only the receiver (who had  $d$ ) can decrypt it

I) Examples

1)  $p = 17$  ,  $q = 11$

2)  $n = p \times q = 17 \times 11 = 187$

3)  $\phi(n) = (p-1)(q-1) = (16 \times 10) = 160$

4)  $e?$   $\gcd(\phi(n), e) = 1$   $1 < e < \phi(n)$   
*should be prime*  $\Rightarrow 1 < e < 160$

options:  $e = 3$  ,  $e = 7$  ,  $e = 11$

5)  $d \cdot e = 1 \pmod{160}$  ;  $d < 160$

$$e \cdot d = 1 + K \cdot \phi(n)$$

$$3 \cdot d = 1 + K \cdot 160 \quad \} \quad K = 1$$

$$3d = 161$$

$$d = 53.66 \text{ (not integer)} \Rightarrow e = 3 \times$$

Try  $e = 7$

$$7 \cdot d = 161$$

$$d = 23 \checkmark \quad \text{Then } e = 7$$

6) Public Key:  $K_U \{e, n\} = \{7, 187\}$

Private Key  $K_R \{d, n\} = \{23, 187\}$

- II) 1)  $p=3$   $q=11$   
 2)  $n=p \times q = 33$   
 3)  $\phi(n) = (p-1)(q-1) = 2 \times 10 = 20$

4)  $e = ?$

$$\left. \begin{array}{l} \gcd(\phi(n), e) = 1 \\ \gcd(20, e) = 1 \\ 1 < e < 20 \end{array} \right\} e \text{ can be } 3, 7, 11, \dots$$

5)  $d \cdot e = 1 \pmod{\phi(n)}$

$$d \cdot e = 1 \pmod{20} \Rightarrow d < 20$$

$$e \cdot d = 1 + K \cdot \phi(n)$$

Let  $K=1$  and  $e=3$

$$3 \cdot d = 1 + 20$$

$$3 \cdot d = 21$$

$$d = 7$$

6) Private Key :  $KR = \{d, n\} = \{7, 20\}$   
 Public Key :  $KU = \{e, n\} = \{3, 20\}$

III) 1)  $p=5$  ,  $q=7$  ,  $M=L$  (12 in alphabetical order)

2)  $n = q \times p = 7 \times 5 = 35$

3)  $\phi(n) = (p-1)(q-1) = 4 \times 6 = 24$

4)  $e = ?$

$$\left. \begin{array}{l} \gcd(24, e) = 1 \\ 1 < e < 24 \end{array} \right\} e \text{ can be } 5, 7, 11, \dots$$

$$5) e \cdot d = 1 \pmod{24} \Rightarrow d < 24$$

$$e \cdot d = 1 + K \phi(n)$$

$$\hookrightarrow \text{let } K=1 \text{ and } e=5$$

$$5 \cdot d = 25$$

$$d = 5 \text{ and } e = 5 \text{ (e and d are the same : not logical)}$$

$$\hookrightarrow \text{let } K=6, e=5$$

$$5 \cdot d = 1 + 6(24)$$

$$5 \cdot d = 145$$

$$d = 29 \checkmark$$

$\Rightarrow$  Public Key:  $K_U = \{5, 35\}$   $\rightarrow$  For encryption  
Private Key:  $K_R = \{29, 35\}$   $\rightarrow$  For decryption

Encryption:

$$C = L^e \pmod{n}$$

$$C = 12^5 \pmod{35}$$

248832	35
248815	7109
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17	

$$\Rightarrow C = 17$$

$$\text{Decryption: } M = C^d \pmod{n} = 17^{29} \pmod{35}$$